

Offset-Symmetric Gaussians for Differential Privacy

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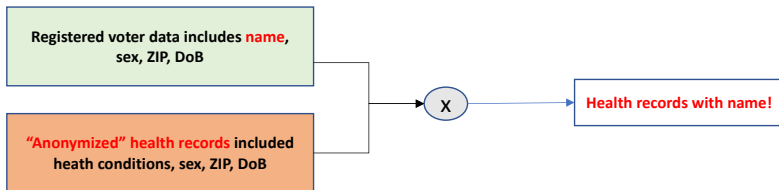
Summary of contributions

- A new “hybrid” differential privacy mechanism is proposed, which is somewhat between Laplace and Gaussian mechanisms.
- It is rooted in the **Gaussian mechanism** → analytical privacy performance derivations.
- It can have **sub-Gaussian tail** → desirable for reducing outliers or post-processing bias.
- At the same **utility** (measured by variance), it has **better** (ϵ, δ) and **better** Rényi differential privacy performance than the Gaussian.
- To achieve the same (ϵ, δ) differential privacy levels, it adds **less** noise to query (measured by noise variance).

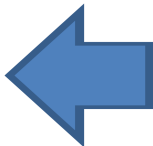
Reconstruction attacks

Facts

- Between 50% to 80% of people in the US are uniquely identified by their full DoB, ZIP code and sex.
- Simple “anonymization” of datasets and publishing them or using them in training algorithms is a privacy risk.
- In mid 1990s, Latanya Sweeney managed to hack into “anonymized” health records of the Governor of MA by linking it with publicly available voter data.



We now know that this publication can be reverse-engineered to reveal the confidential database.



	Count	Median	Mean
Total	7	30	38
# Female	4	30	33.5
# male	3	30	44
# black	4	51	48.5
# white	3	24	24
Married	4	51	54
Black F	3	36	36.7

This table can be expressed by 164 equations.
Solving those equations takes
0.2 seconds on a 2013 MacBook Pro.

Facts

- Reconstructed over 300 million records from the 2010 census publicly released data
- Used 4 commercial databases which included real people's name, address, age, sex
- Linked reconstructed records with commercial databases including name, address, age, sex, ethnicity and race
- Comparing with confidential US census data could get all variables (including race & ethnicity) right for 17% of the US population

Source: <https://simson.net/ref/2019/2019-07-16%20Deploying%20Differential%20Privacy%20for%20the%202020%20Census.pdf>

Deploying Differential Privacy for the 2020 Census of Population and Housing

Simson L. Garfinkel
Senior Scientist, Confidentiality and Data Access
U.S. Census Bureau

July 16, 2019
Privacy Enhancing Technologies Symposium
Stockholm, Sweden 2019

US Census in 2020

- “Every person matters for federal funding.”
- To preserve privacy: “Imaginary people will be added to some locations and real people will be removed from others.”
- “Minorities and rural areas at most risk.”

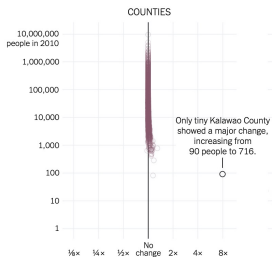


Source: <https://www.nytimes.com/interactive/2020/02/06/opinion/census-algorithm-privacy.html#commentsContainer>

Outliers Matter

- Y axis: true population count (log scale)
- X axis: reported count relative to the true population count (log scale)

Large populations mostly stayed the same size under the privacy algorithm.

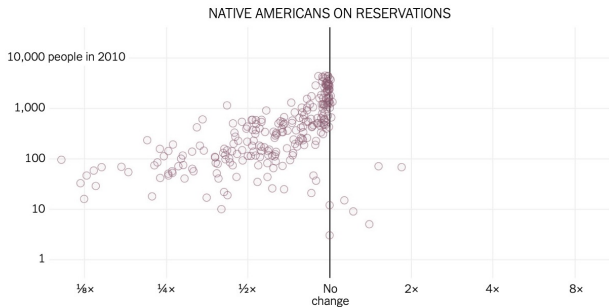


Small populations were affected more by the algorithm. Some doubled in size; others were halved or even disappeared.



Source: <https://www.nytimes.com/interactive/2020/02/06/opinion/census-algorithm-privacy.html#commentsContainer>

Under the privacy algorithm, almost all small- and medium-size reservations showed **fewer Native American inhabitants**.



Source: <https://www.nytimes.com/interactive/2020/02/06/opinion/census-algorithm-privacy.html#commentsContainer>

ABS Perturbation Methodology Through the Lens of Differential Privacy

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Abstract. The Australian Bureau of Statistics (ABS), like other national statistical offices, is considering the opportunities of differential privacy (DP). This research considers the Australian Bureau of Statistics (ABS) TableBuilder perturbation methodology in a DP framework. DP and the ABS perturbation methodology are applying the same idea – infusing noise to the underlying microdata – to protect aggregate statistical outputs. This research describes some differences between these approaches. Our findings show that noise infusion protects against disclosure risks in the aggregate Census Tables. We highlight areas of future ABS research on this topic.

1 Introduction

The world is witnessing an explosion in the automated collection of personal data; a reduction in the cost of high-powered computational resources; and the increased frequency and sophistication of data attacks. It is a regular occurrence for cyber attacks to make the news. Naturally, this elevates the public concern over privacy and how personal information is used once collected. Public trust in the Australian Bureau of Statistics (ABS) to protect the data it collects from providers, is a cornerstone to the ABS mission. The US Census Bureau (USCIB) recently announced – via its Scientific Advisory Committee – that it would protect the publications of the 2018 End-to-End Census Test (E2E) using differential privacy (DP). The E2E test is a dress rehearsal for the 2020 Census¹ (Abowd, 2018, p.2). In light of this announcement, many National Statistical Offices (NSOs), including the Office for National Statistics and Statistics New Zealand, are investigating DP approaches to protecting their Census outputs.

The ABS has been exploring the possibility of adopting a DP approach to improve existing confidentiality methodologies, particularly those that apply to the ABS TableBuilder. This paper contributes to the current discussion on DP by providing the current ABS perspective.

How to educate researchers, the public and politicians about such tradeoff?

Managing the Tradeoff



Basic notation

- x is a database, counting query $q(x)$ returns true count n :

$$\underbrace{M(x)}_{\text{mechanism noisy output}} = \underbrace{n}_{q(x) \text{ is true count}} + \underbrace{Y}_{\text{noise}}$$

- Data-independent noise added to the true count.

$$\Pr(M(x) = i | n) = \Pr(Y = i - n | n) \underbrace{=}_{\text{independence}} \Pr(Y = i - n)$$

- For every two neighbouring datasets $x \sim x'$ differing on the attribute of one individual (e.g., **smokes** or does not smoke), the output of the mechanism should be almost indistinguishable.
- ϵ -DP:

$$e^{-\epsilon} \leq \frac{\Pr(M(x) = i)}{\Pr(M(x') = i)} = \frac{\Pr(M = i | n)}{\Pr(M = i | n + 1)} \leq e^{\epsilon}, \forall i, n.$$

Formal definition

- A randomized mechanism is $M : \mathcal{X}^n \rightarrow \mathcal{Y}$ (n elements in dataset).
- If for all neighboring datasets $x \sim x' \in \mathcal{X}^n$ and all events $E \subset \mathcal{Y}$, we have

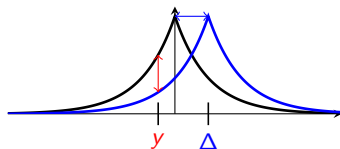
$$\mathbb{P}[M(x) \in E] \leq e^\epsilon \mathbb{P}[M(x') \in E] + \delta,$$

then we say M satisfies (ϵ, δ) -DP.

- **Pure** differential privacy: $(\epsilon, 0)$ -DP.
- **Approximate** differential privacy: (ϵ, δ) -DP, $0 < \delta < 1$.

Basic notation

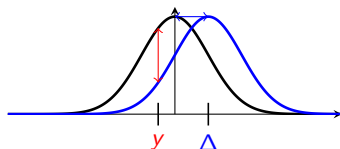
- Assume two neighboring datasets $x \sim x'$ differing on one element.
- Query function q is called; Laplace noise $Y \leftarrow \mathcal{L}(0, \lambda)$ is added to q .
- $q(x) = 0 \quad \Rightarrow \quad M(x) = 0 + Y.$
- $q(x') = \Delta \quad \Rightarrow \quad M(x') = \Delta + Y.$



- The **worst-case ratio** of the two probability density functions will determine the differential privacy (DP) performance.
- $\lambda = \frac{\Delta}{\epsilon} \quad \Rightarrow \quad \epsilon\text{-DP}.$

Basic notation

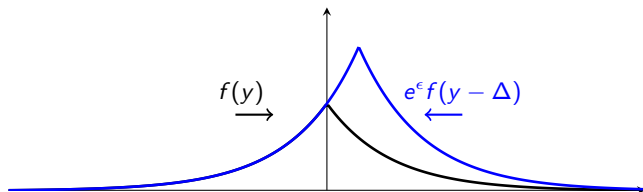
- $M(x) = q(x) + Y$, $Y \leftarrow \mathcal{N}(0, \sigma_{\text{gauss}}^2)$.



- There is **no finite** ϵ for which the ratio of the two probability density functions is bounded by $e^\epsilon \Rightarrow$ We need to use either (ϵ, δ) -DP or Rényi DP frameworks. In other words:

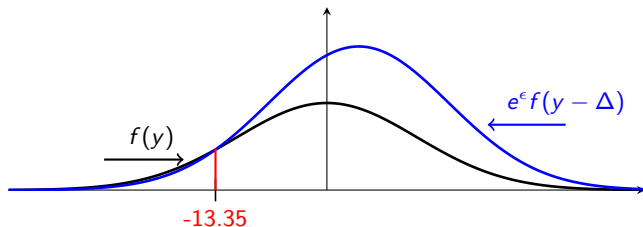
$$E = \{y : \frac{f(y)}{f(y - \Delta)} > e^\epsilon\} \neq \emptyset$$

Pros and cons of Laplace DP mechanism



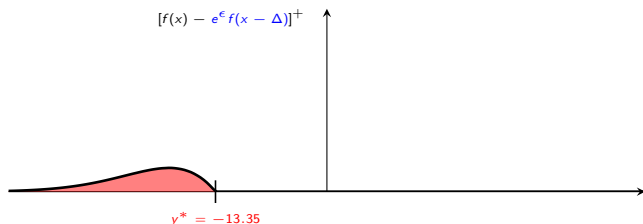
- **Pro:** Can achieve pure differential privacy: $(\epsilon, 0)$ -DP:
The blue curve never goes below the black curve.
- Easy-to-design: e.g., $\epsilon = 0.5$, $\Delta = 1$, Laplace scale $\lambda = \frac{\Delta}{\epsilon} = 2$.
- Utility optimal in some scenarios.
- **Con:** Heavy tail \rightarrow outliers and bias \rightarrow challenges in post processing.

Pros and cons of Gaussian DP mechanism



- **Pro:** Sub-Gaussian tail \rightarrow less outliers or bias after post-processing.
- **Con:** Only approximate (ϵ, δ) -DP is possible \rightarrow δ is a measure of impossibility of pure ϵ -DP.
- For $\epsilon = 0.5$, $\Delta = 1$ and $\sigma_{\text{gauss}}^2 \approx 27.7$, values $y < -13.35$ contribute to $\delta \approx 3.2 \times 10^{-4}$ (figure not to scale).

Computation of δ



- Critical value of y , below which all values contribute to δ :

$$y^* = \frac{\Delta}{2} - \frac{\epsilon \sigma_{\text{gauss}}^2}{\Delta} = \frac{1}{2} - \frac{0.5 * 27.7}{1} = -13.35$$

- δ is given by the **area under the curve**:

$$\delta(\epsilon) = \int_{-\infty}^{y^*} (f(x) - e^\epsilon f(x - \Delta)) dx,$$

where f denotes the pdf of the Gaussian $\mathcal{N}(0, \sigma_{\text{gauss}}^2)$.

- **Figure not to scale.**

- Discrete Gaussian is used for the release of 2020 US Census.
- One main reason for switching from Laplace to Gaussian was the “heavy tail” problem of Laplace leading to outliers and bias during processing.
- δ should be much smaller than $1/n$, where n is the size of population.

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Research Question: Can we keep light tail of Gaussian but improve on its DP-variance tuple (joint) performance?

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Research Question: Can we keep light tail of Gaussian but improve on its DP-variance tuple (joint) performance?

The answer is YES!

Definition and properties

$$f(y) = \underbrace{\frac{1}{S'}}_{\text{normalization}} \underbrace{e^{-\frac{y^2}{2\sigma^2}}}_{\text{Gaussian}} \underbrace{e^{-\frac{|y|m}{\sigma^2}}}_{\text{Laplace factor } \mathcal{L}(0, \frac{m}{\sigma^2})} \quad (1)$$

- Point-wise product of a Gaussian pdf and a Laplace pdf.
- Original idea: moderate or slow down the Gaussian decay to reduce δ through factoring in a Laplace pdf.
- Two parameters in pdf: σ^2 and m (more DoF in privacy-utility tradeoff).
- Can be interpreted in different ways as we will see next (including why it is called Offset Symmetric Gaussian Tails (OSGT)).

Alternative Description of pdf: The Reason behind OSGT Naming - 1

Towards the OSGT pdf

Our original definition can be rewritten as:

$$f(y) = \frac{1}{S'} e^{-\frac{y^2}{2\sigma^2} - \frac{|y|m}{\sigma^2}} \quad \text{Expand } |y| \text{ cases} \Rightarrow$$

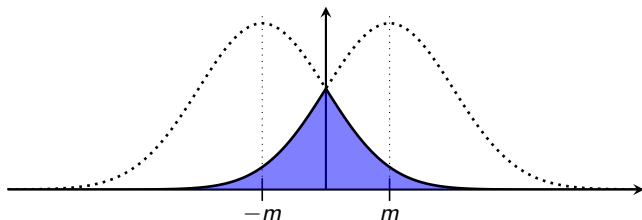
$$= \begin{cases} \frac{1}{S'} e^{-\frac{y^2}{2\sigma^2} + \frac{ym}{\sigma^2}}, & y \leq 0, \\ \frac{1}{S'} e^{-\frac{y^2}{2\sigma^2} - \frac{ym}{\sigma^2}}, & y > 0, \end{cases} \quad \text{Multiply by 1 or by } e^{+\frac{m^2}{2\sigma^2}} e^{-\frac{m^2}{2\sigma^2}}$$

$$= \begin{cases} \frac{e^{+\frac{m^2}{2\sigma^2}}}{S'} e^{-\frac{(y-m)^2}{2\sigma^2}}, & y \leq 0, \\ \underbrace{\frac{e^{+\frac{m^2}{2\sigma^2}}}{S'}}_{\text{New constant } 1/S} e^{-\frac{(y+m)^2}{2\sigma^2}}, & y > 0, \end{cases}$$

OSGT pdf

$$f(y) = \begin{cases} \frac{1}{S} e^{-\frac{(y-m)^2}{2\sigma^2}}, & y \leq 0, \\ \frac{1}{S} e^{-\frac{(y+m)^2}{2\sigma^2}}, & y > 0, \end{cases}$$

Normalization $S = \sqrt{2\pi\sigma^2} 2Q(m/\sigma)$ takes care of the area under the Gaussian tails. Note the resulting distribution has **zero mean**.



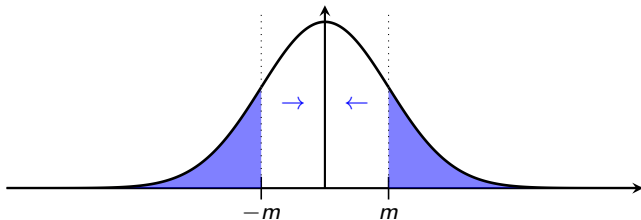
The Gaussian Q -function:

$$Q(x) = \frac{1}{\sqrt{2\pi}} \int_x^\infty e^{-t^2/2} dt$$

Sampling: rejecting the non-tail values of $\mathcal{N}(0, \sigma^2)$

Rejection, followed by a shift

- This view will be useful in **sampling from the Gaussian**, rejecting values between $[-m, m)$, and **shifting the acceptable tail values** to obtain a sample from the OSGT.



Approximation at small $|y| \approx$ scaled Laplace

$$\begin{aligned}
 f(y) &= \underbrace{\frac{1}{S'}}_{\text{normalization}} \underbrace{e^{-\frac{y^2}{2\sigma^2}}}_{\text{Gaussian}} \underbrace{e^{-\frac{|y|m}{\sigma^2}}}_{\text{Laplace}} \\
 &= \frac{1}{S'} e^{-\frac{|y|(|y|+2m)}{2\sigma^2}} \\
 &\approx \frac{1}{S'} e^{-\frac{|y|(\cancel{|y|}+2m)}{2\sigma^2}}, \quad |y| \ll 2m \\
 &= \frac{1}{S'} e^{-\frac{m|y|}{\sigma^2}}, \quad |y| \ll 2m.
 \end{aligned}$$

The shape of the proposed pdf

- All three distributions are compared at the same variance.
- The “Laplace-like” behaviour of the **new distribution** at small $|y|$ can be observed.

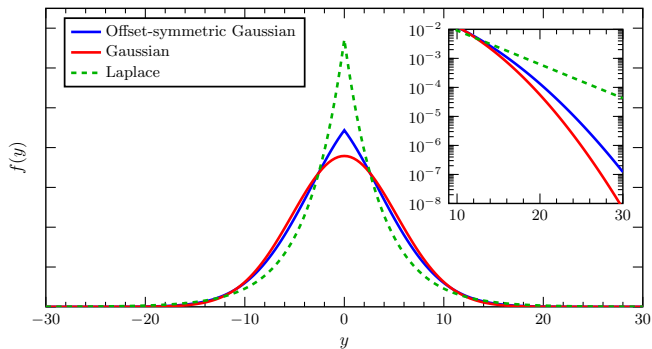


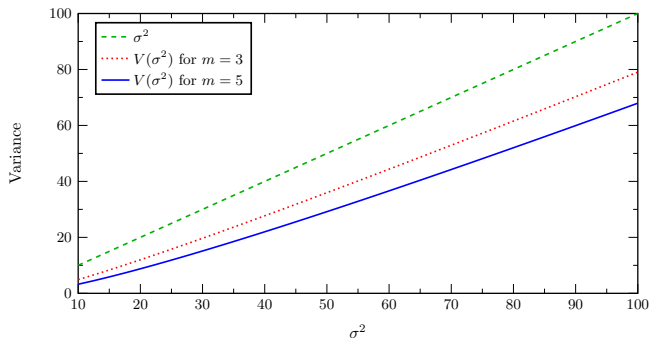
Figure: $\sigma^2 = 40$, $m = 3$, but **actual variance is $V \approx 27.7$** .

Variance of OSGT random variable

Variance of OSGT is smaller than its input parameter σ^2

Let $Y \leftarrow \mathcal{T}(0, V(m, \sigma^2))$ be a zero-mean OSGT distributed random variable. Its variance is given by

$$\mathbb{E}[Y^2] = V(m, \sigma^2) = \sigma^2 + m^2 - \frac{m\sigma \exp\left(-\frac{m^2}{2\sigma^2}\right)}{\sqrt{2\pi}Q\left(\frac{m}{\sigma}\right)} < \sigma^2. \quad (2)$$



Proposition

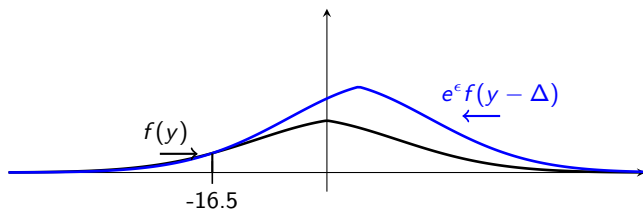
- The OSGT distribution $\mathcal{T}(0, V(m, \sigma^2))$ is σ^2 -sub-Gaussian if

$$\frac{m}{\sigma} \leq Q^{-1}(0.25) \approx 0.6745.$$

- In this case, we have:

$$\mathbb{P}_{Y \leftarrow \mathcal{T}(0, V(m, \sigma^2))} [Y \geq y] \leq \exp\left(-\frac{y^2}{2\sigma^2}\right).$$

Computing δ for the OSGT mechanism - 1



- At $\epsilon = 0.5$, $\Delta = 1$, $m = 3$, and $\sigma^2 \approx 40$, values $y < -16.5$ contribute to $\delta \approx 6.8 \times 10^{-5}$.
- The actual variance is $V = 27.7$.
- Recall, for the same variance $\sigma_{\text{gauss}}^2 = 27.7$, same $\epsilon = 0.5$ and $\Delta = 1$, the Gaussian gives $y < -13.35$ and $\delta \approx 3.2 \times 10^{-4}$.

Numerical evaluation of (ε, δ) for OSGT versus Gaussian

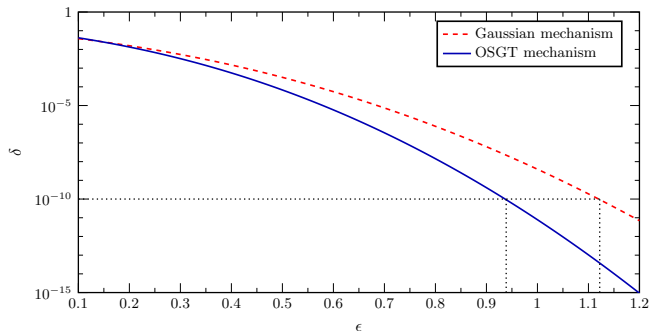


Figure: $\Delta = 1$, $m = 3$, $\sigma^2 = 40$, $V(m, \sigma^2) = \sigma_{\text{gauss}}^2 \approx 27.7$.

- The OSGT mechanism achieves $(\varepsilon, \delta) \approx (0.94, 10^{-10})$.
- The Gaussian mechanism achieves $(\varepsilon, \delta) \approx (1.12, 10^{-10})$.

Numerical evaluation of (ϵ, δ) for OSGT versus Gaussian - 2

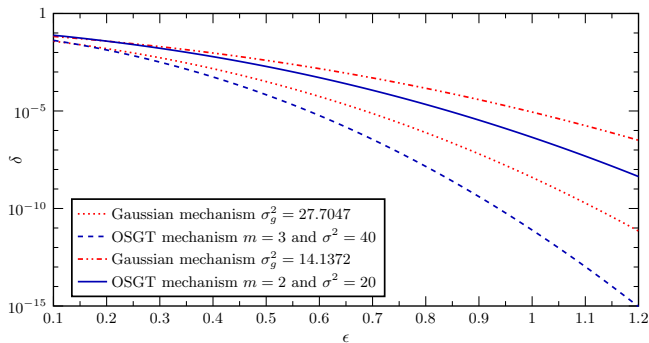


Figure: (ϵ, δ) -DP performance of OSGT and Gaussian mechanisms at the same variance.

- The OSGT mechanism achieves better (ϵ, δ) at the same variance.

Numerical evaluation of (ϵ, δ) for OSGT versus Gaussian - 3

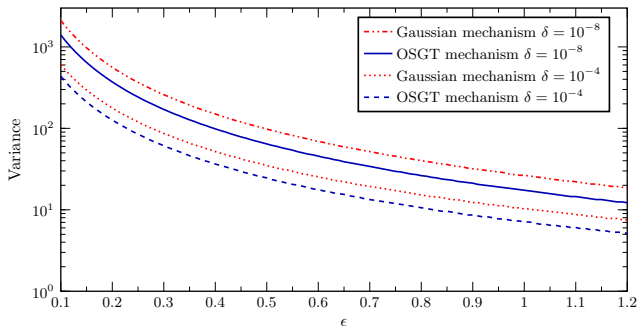
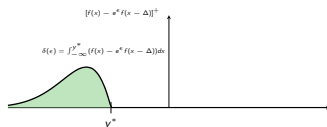


Figure: Variance performance of OSGT and Gaussian mechanisms at the same (ϵ, δ) -DP level.

- The OSGT mechanism needs smaller variance (adds less noise; proxy for utility) at the same (ϵ, δ) .

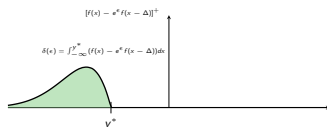
Computing δ for the OSGT Mechanism



Analytical formula for δ , single-variate OSGT mechanism (1/2 cases shown)

$$\delta^T(\epsilon) = \frac{1}{2Q\left(\frac{m}{\sigma}\right)} \left(\underbrace{Q\left(\frac{\sigma\epsilon}{\Delta} - \frac{\Delta}{2\sigma}\right) - e^{\epsilon} Q\left(\frac{\sigma\epsilon}{\Delta} + \frac{\Delta}{2\sigma}\right)}_{\text{decreasing function of } \sigma} \right), \quad \epsilon > \frac{\Delta^2}{2\sigma^2} + \frac{\Delta m}{\sigma^2}.$$

Computing δ for the OSGT Mechanism



Analytical formula for δ , single-variate OSGT mechanism (1/2 cases shown)

$$\delta^{\mathcal{T}}(\epsilon) = \frac{1}{2Q\left(\frac{m}{\sigma}\right)} \left(\underbrace{Q\left(\frac{\sigma\epsilon}{\Delta} - \frac{\Delta}{2\sigma}\right) - e^{\epsilon} Q\left(\frac{\sigma\epsilon}{\Delta} + \frac{\Delta}{2\sigma}\right)}_{\text{decreasing function of } \sigma} \right), \quad \epsilon > \frac{\Delta^2}{2\sigma^2} + \frac{\Delta m}{\sigma^2}.$$

Recall that for a fair utility comparison: $\sigma_{\text{gauss}}^2 = V(m, \sigma^2) < \sigma^2$.

Comparison with δ of the Gaussian mechanism

$$\delta^{\mathcal{N}}(\epsilon) = Q\left(\frac{\sigma_{\text{gauss}}\epsilon}{\Delta} - \frac{\Delta}{2\sigma_{\text{gauss}}}\right) - e^{\epsilon} Q\left(\frac{\sigma_{\text{gauss}}\epsilon}{\Delta} + \frac{\Delta}{2\sigma_{\text{gauss}}}\right), \quad \epsilon \geq \frac{\Delta^2}{2\sigma_{\text{gauss}}^2}.$$

Definition

- The k -dimensional query function $q : \mathcal{X}^n \rightarrow \mathbb{R}^k$.
- The OSGT mechanism is obtained as $M(x) = q(x) + Y$.
- Each noise component Y_i in $Y = (Y_1, \dots, Y_k)$ is drawn from an i.i.d. scalar OSGT distribution $\mathcal{T}(0, V(m, \sigma^2))$, leading to:

$$f_{M(x)}^{\mathcal{T}}(y) = \frac{1}{S'^k} \exp \left(\underbrace{-\frac{\|y - q\|_2^2}{2\sigma^2}}_{k\text{-dim Gaussian}} \underbrace{-\frac{m \|y - q\|_1}{\sigma^2}}_{k\text{-dim Laplace}} \right).$$

Multidimensional OSGT mechanism

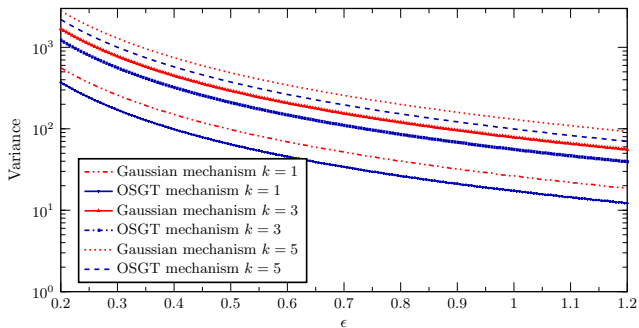


Figure: Variance performance of OSGT and Gaussian mechanisms at the same (ϵ, δ) -

- The OSGT mechanism needs smaller variance (adds less noise; proxy for utility) at the same better (ϵ, δ) for k -dimensional query.
- However, the improvement seems to diminish for large k .
- The method for computation of δ is numerical.

Single-letter RDP upper bound

The Rényi DP of the OSGT mechanism is upper bounded by

$$\underbrace{D_\alpha(M(x) \| M(x'))}_{\text{Rényi DP}} \leq \alpha \frac{\Delta_2^2}{2\sigma^2} + \zeta, \quad \zeta = \frac{k}{\alpha - 1} \log \left(\frac{1 - Q\left(\frac{m}{\sigma}\right)}{Q\left(\frac{m}{\sigma}\right)} \right).$$

where $D_\alpha(M(x) \| M(x'))$ is the Rényi divergence of order α .

Gaussian achieves RDP as $\alpha \frac{\Delta_2^2}{2\sigma_{\text{gauss}}^2} > \alpha \frac{\Delta_2^2}{2\sigma^2}$.

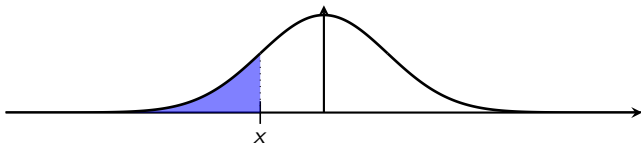
Key equations

- The Rényi Divergence for k -dimensional OSGT mechanism can be analytically written as

$$D_\alpha(M(x) \| M(x')) \leq \frac{\alpha \Delta^2}{2\sigma^2} + \frac{k}{\alpha - 1} \log \left(\frac{\sqrt{2\pi\sigma^2} \bar{B}}{S} \right)$$

- where:

$$\begin{aligned} \bar{B} = & \Phi((-m + (\alpha - 1)\Delta)/\sigma) + \Phi((-m - \alpha\Delta)/\sigma) + \\ & e^{\alpha(\alpha-1)(4m\Delta+4m^2)/2\sigma^2} \times \left(\Phi\left(\frac{\alpha\Delta - m(1-2\alpha)}{\sigma}\right) - \Phi\left(\frac{(\alpha-1)\Delta - m(1-2\alpha)}{\sigma}\right) \right). \end{aligned}$$



Φ is the Gaussian cumulative distribution function (cdf):

$$\Phi(x) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-t^2/2} dt = 1 - Q(x).$$

Components of Rényi divergence for OSGT mechanism

Key equation

$$D_\alpha(M(x) \| M(x')) \leq \underbrace{\alpha \frac{\Delta_2^2}{2\sigma^2}}_{\alpha\rho} + \underbrace{\frac{k}{\alpha-1} \log\left(\frac{\sqrt{2\pi\sigma^2} \bar{B}}{S}\right)}_{\zeta}$$

We numerically evaluate the scale of ζ and $\alpha\rho$ as a function of α .

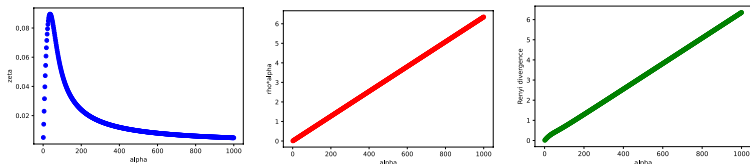


Figure: $k = 8, m = 15, \sigma^2 = 630, \Delta = 1, \Delta_2^2 = k\Delta = 8$. It can be observed the effect of ζ in the overall Rényi divergence is rather small.

Analytical computation of Rényi DP to get a bound on δ

$$D_\alpha(M(x)||M(x')) \leq \tau \quad \Rightarrow \quad \delta(\varepsilon) = \frac{\exp((\alpha - 1)(\tau - \varepsilon))}{\alpha - 1} \left(1 - \frac{1}{\alpha}\right)^\alpha$$
$$\tau = \frac{\alpha \Delta_2^2}{2\sigma^2} + \frac{k}{\alpha - 1} \log \left(\frac{\sqrt{2\pi\sigma^2} \bar{B}}{S} \right).$$

For more details and applications (such as composition and noise filtering), please see our paper published in 2022 in IEEE Trans. Information Forensics and Security (TIFS) with the same title.